

- 1 -

DATA FLOW BETWEEN A COMMUNICATION NODE AND A MOBILE NODE
IN A MOBILE NETWORK

Field of the Invention

5

This invention relates to telecommunication systems in which data flows between mobile nodes, for example personal digital assistants with wireless communication capability, and a data network, for example the Internet.

10

Background of the Invention

The Internet is becoming more and more popular, and users increasingly wish to access the Internet whilst on the move. Different types of mobile nodes (i.e. mobile communication units) may be employed for this purpose, for example a mobile telephone or a personal digital assistant (PDA) with wireless communication capability.

20 Increasingly, mobile users are accessing the Internet via different types of fixed or wireless access networks, for example a cellular radio communication network, such as a Universal Mobile Telecommunication System (UMTS) network, a HiperLAN/2 or IEEE 802.11b local area network, a
25 Bluetooth local communication system, or fixed accesses such as the Ethernet, and so on. The data route between the mobile node and the Internet further comprises an Internet protocol subnet (IP subnet), such that the route is as follows: mobile node - access network - IP subnet -
30 Internet (and reverse order for the data route from the Internet to the mobile node).

- 2 -

It is currently possible to seamlessly handover accessing of the Internet from one access network to another, for example by using a protocol known as Mobile-IP.

5 Traditional mobility support aims to provide continuous Internet connectivity to mobile hosts, thereby allowing individual mobile users to connect to the Internet whilst being mobile and moving their Internet access location. In contrast, network mobility support is concerned with
10 situations where an entire network changes its point of attachment to the Internet topology and thus its accessibility in the topology. Such a network in movement can be called a Mobile Network.

15 There exist a large number of scenarios where such Mobile Networks exist. For example, a Personal Area Network (PAN, i.e. a network of several personal devices attached to an individual) is known whereby the PAN changes its point of attachment to the Internet topology whilst the
20 user is walking in a shopping mall. In addition, a network may be embedded in a bus or aircraft, providing on-board Internet access to passengers. A passenger may use a single communication device (e.g. a laptop) or be itself a Mobile Network (e.g. a PAN). Notably, this
25 configuration illustrates a case of a Mobile Network visiting a Mobile Network (i.e. nested mobility).

As such, a Mobile Network can be defined as a set of nodes composed of one or more IP-subnets attached to a
30 Mobile Router (MR). These IP subnets may also be viewed as a mobile unit, with respect to the rest of the

- 3 -

Internet, i.e. a MR and all its attached nodes (so called Mobile Network Nodes or MNNs).

A document [1] authored by Thierry Ernst, Hong-Yon Lach,
5 IETF Internet-Draft draft-ernst-monet-terminology-00.txt,
February 2002, describes a list of definitions for the Mobile Network terminology that can be applied in this application. In particular, the following terms may be defined as follows:

10

(i) A Local Fixed Node (LFN):

A node permanently located within the Mobile Network and that does not change its point of attachment. A LFN can either be a Local Fixed Host (LFH) or a Local Fixed
15 Router (LFR).

(ii) A Local Mobile Node (LMN):

A local mobile node is one that belongs to the Mobile Network and changes its point of attachment from a link
20 within the Mobile Network to another link within or outside the Mobile Network. In this regard, it can be assumed that the home link of the LMN is a link within the Mobile Network. A LMN can either be a Local Mobile Host (LMH) or a Local Mobile Router (LMR).

25

(iii) A Visiting Mobile Node (VMN):

A VMN is one that does not belong to the Mobile Network, and changes its point of attachment from a link outside the Mobile Network to a link within the Mobile Network
30 (i.e. the home link of the VMN is not a link within the Mobile Network). A VMN that attaches to a link within the Mobile Network obtains an address on that link. A

- 4 -

VMN can either be a Visiting Mobile Host (VMH) or a Visiting Mobile Router (VMR).

(iv) A Mobile Network Prefix:

5 A bit string that consists of a number of initial bits of an IP address, which identifies a Mobile Network within the Internet topology. Nodes belonging to the Mobile Network (i.e. at least MR, LFNs and LMNs) share the same IPv6 "network identifier". For a single mobile IP-
10 subnet, the Mobile Network Prefix is the "network identifier" of this subnet.

(v) An Egress Interface of a MR:

This is the interface attached to the home link if the
15 Mobile Network is at home. Alternatively, it is the interface attached to a foreign link if the Mobile Network is in a foreign network.

(vi) An Ingress Interface of a MR:

20 This is the interface attached to a link inside the Mobile Network. The concept of egress and ingress interfaces can be extended to other types of nodes within a mobile network as follows:

(a) All network interfaces of hosts (fixed or
25 mobile) in a mobile network are said to be egress interfaces. None of them is said to be an ingress interface.

(b) A fixed router in a mobile network (i.e. LFR) may have several egress and ingress interfaces. The
30 type of each interface should be pre-configured on such routers. An egress interface is usually an interface on the shortest path to a Mobile Router. A fixed router in

- 5 -

a mobile network should have at least one egress interface.

(vii) A MR may have multiple Egress and Ingress
5 interfaces.

(viii) A mobile network may have multiple MRs.

Recently, there has been a lot of interest and research
10 into the Mobile IPv6 specification, as described in the document [2] authored by David B. Johnson: IETF Internet-Draft draft-ietf-mobileip-ipv6-15.txt, July 2001. A major concern with Mobile IPv6 is that the research has proven that the IPv6 standard is currently unable to
15 adequately address network mobility. In particular, the document [3] authored by Thierry Ernst and Hong-Yon Lach: IETF Internet-Draft draft-ernst-mobileip-v6-network-02.txt, June 2001, details problems encountered with Mobile-IPv6 in supporting Mobile Networks.

20 In summary, it has been determined that even if a MR's Home Agent (HA) is able to intercept packets addressed to MNs that are operating behind the MR, the MR's HA is clearly unable to encapsulate them to the care-of-address
25 of the appropriate MR. Note that every data packet has a source address and a destination address. A tunnelled packet is a packet that encapsulates in it another packet. Thus, the encapsulating packet has a pair of source and destination addresses. A further encapsulated
30 packet has additional source and destination addresses.

- 6 -

The lack of knowledge of a true location of a particular MNN results from the HA not knowing any (never mind a preferred) data route to the Mobile Network, once it has moved to a 'visited' location. Unfortunately, when a MR registers with its HA the MR only informs the HA to record a host-specific route in its routing table. The inventors of the present invention have recognised that a preferred network route generated using the mobile network address (prefix, prefix length and care-of-address) of the appropriate MR would greatly assist in this matter.

In the field of this invention, the Mobile IPv4 specification, detailed in [4] C. Perkins, IETF RFC 3220, "IP Mobility Support for IPv4", Standards Track, January 2002, describes how Network Mobility can be supported in the case of IPv4 mobility. However, it has also been determined that IPv4 does not support route optimisation for MNNs behind the MR.

Thus, all the (incoming and outgoing) traffic between a MNN and its corresponding nodes (CNs) is passed to the MR's Home Agent. This problem is exacerbated in the case of nested mobility, which is where a CN wishes to pass data to a MNN that is behind a number of MR links, or a mobile node visiting a mobile network. In the case of nested mobility, the packet will thus be encapsulated several times as a result of being routed through all of the Home Agents of all the nested MRs. This is clearly inefficient routing.

- 7 -

Similarly, in the case of IPv6, the document [5] authored by T. J. Kniveton: IETF Internet-Draft draft-kniveton-mobrrtr-00.txt, November 2001 describes how to support mobile networks with no modifications to Mobile IP (v4 or v6). The mechanism proposed to address the support of nested mobility is described with respect to FIG. 2. It suffers from the same problem of multiple tunnelling by the Home Agents of the nested mobile routers as described in the previous paragraph. When a MR, for example MR2 260, has attached to a visited network 110 (via another Mobile Network MR1 link), a bi-directional tunnel 210, 215, 220 is established between MR2 260 and its HA - HA2 250. When a node, for example LFN2 165, is attached to MR2 260 and wishes to send an IP packet to a CN, say CN2 255, via the Internet 115, that packet is tunnelled by MR2 260 (to HA2 250) and again by MR1 150 (to HA1 240). The multiple-tunnelled data packet is then passed to the HA of the latest MR to tunnel the data, namely to HA1 240. HA1 240 then forwards it to the intended recipient CN2 255 via the source MR's HA, namely HA2 250.

This proposed data routing method, and the problems associated with it, are best described by way of an example. Thus, let us assume that LFN2 165 sends a data packet to CN2 255. The data packet is first routed towards MR2 260. The data packet is then tunnelled by MR2 260 to be sent to HA2 250. This tunnelling process of the data packet from MR2 260 to HA2 250 is itself, by necessity, first routed 210 towards its linked MR, namely MR1 150. MR1 150 further tunnels the data and forwards 215 the multiple-tunnelled packet to its HA, namely HA1 240. HA1 240 de-tunnels the data packet, as tunnelled by

- 8 -

MR1 150 and forwards 220 the partially de-tunnelled
packet to its original intended recipient, HA2 250. HA2
250 further de-tunnels the packet, as tunnelled by MR2
260, and forwards 225 the wholly de-tunnelled data packet
5 to CN2 255.

Clearly, the solution proposed does not provide any
support for route optimisation, since both inbound and
outbound packets are routed through the Home Agent of
10 both MRs 150, 260. In fact, packets from CN2 255
addressed to LFN2 165 will follow the same path (in
reverse order) and will then be encapsulated by each Home
Agent 240, 250 of each of the nested MRs 150, 260. Once
again, this is clearly inefficient routing, particularly
15 for a practical situation whereby there may be many more
than two levels in the nested network.

The solution presented in a document authored by Thierry
Ernst and Hong-Yon Lach: IETF Internet-Draft draft-ernst-
20 mobileip-v6-network-02.txt, June 2001, proposes a means
for supporting network mobility in the framework of
Mobile IPv6. This solution introduces the following
concept: when a Mobile Router roams to a visited network,
it sends a Prefix Scope Binding Update to its Home Agent
25 (HA). Unlike a classical Mobile IPv6 Binding Update
message, a Prefix Scope Binding Update does not bind a
home address with a care-of address.

In contrast, the MR Prefix is bound with the MR care-of
30 address, for a particular MR. Upon reception of a packet
whose prefix matches with the MR prefix, the Home Agent
must then tunnel the packet to the MR care-of-address

- 9 -

that has been identified as being able to deliver the
tunnelled packet to the intended recipient. Likewise, a
MR may send prefix-scope BUs to the corresponding nodes
of the nodes it serves. This solution brings a more
5 efficient support of mobile networks to Mobile IPv6,
since it may provide a limited improvement to route
optimisation.

However, the scope of this draft has explicitly excluded
10 the case of nested mobility, which presents a significant
hurdle to efficient route optimisation. As such, the
solution proposed in the Thierry Ernst document, IETF
Internet-Draft draft-ernst-mobileip-v6-network-02.txt,
June 2001, fails to provide a useful solution to route
15 optimisation in many practical situations. Again, the
problems that emanate from this document, particularly in
the case of nested mobility, are best highlighted in an
example situation.

20 Referring now to FIG. 3, a mechanism for routing data
packets in an IPv6 network using the proposal of Thierry
Ernst: IETF Internet-Draft draft-ernst-mobileip-v6-
network-02.txt, June 2001. Notably, the problems
emanating from using this mechanism in a nested mobility
25 are highlighted.

A mobile router MR1 150 is attached to a visited link
110. A mobile router MR2 260 is attached to MR1's link
155. A local fixed node LFN2 165 is attached to MR2's
30 link 230. Again, let us assume that LFN2 165 is
attempting to communicate with a corresponding node CN2
255.

- 10 -

Let us further assume that, at the beginning of the simple scenario detailed in FIG. 3, MR1 150 and MR2 260 have already sent BU messages to their respective HAS
5 240, 250. That is, HA1 240 knows that MR1's 150 prefix is reachable at MR1's care-of address. Similarly, HA2 250 knows that MR2's 260 prefix is reachable at MR2's care-of address.

10 When a data packet is sent from CN2 255 to LFN2 165, CN2 255 has no knowledge about LFN2's 165 actual location. Thus, the data packet that it sends is therefore routed 325 towards home link-2 105. HA2 250 intercepts the data packet and tunnels it to MR2's care-of address. This can
15 be understood, as HA2 250 knows that MR2's prefix is reachable at MR2's care-of address.

This tunnelled packet (from HA2 250 to MR2's 260 care-of address) is routed 320 toward link-1 245, since MR2's 260
20 care-of address matches MR1's 150 prefix. HA1 240 intercepts the data packet and tunnels it to MR1's care-of address, namely towards the visited link 110, since HA1 240 knows that MR1's prefix is reachable at MR1's care-of address.

25 MR1 150 then de-tunnels the data packet received from HA1 240. MR1 150 then forwards the content to the original recipient, MR2 260. Meanwhile, MR1 150 sends a Binding Update to the sender of the encapsulated packet (that is
30 HA2 250) to inform it that MR1's prefix is reachable at MR1's care-of address. Note that from MR1's point of view, HA2 250 is a Correspondent Node and not its Home

- 11 -

Agent (i.e. the 'H' bit in the BU is not set). It is up to the HA2 250 to accept this Binding Update or not.

MR2 260 de-tunnels the data packet that it received (i.e. the portion of the data packet that had been encapsulated by HA2 250) and forwards the content (the initial packet from CN2 255) to LFN2 165. Meanwhile, MR2 260 sends a Binding Update message to the sender of the encapsulated packet (that is, CN2 255) to inform it that MR2's prefix (covering the LFN2 165 address) is reachable at MR2's care-of address. This information is stored in the CN's binding cache 370.

Once an initial packet has reached its destination, transmission of a second or subsequent packet from CN2 255 to LFN2 165 leads to the scenario depicted in FIG. 4. After having reviewed its binding cache 370, CN2 255 recognises that LFN2 165 is reachable at MR2's care-of address.

Thus, it sends the data packet to MR2's care-of-address with a routing header for LFN2 165. MR2's care-of-address belongs to MR1's link and is therefore routed towards Home Link-1 245. In this manner, a minor improvement to route optimisation is achieved by the bypassing of the transmission of the data packet to, and from, HA2 250.

HA1 240 then intercepts the data packet and tunnels the packet to MR1's care-of address, since HA1 240 knows that MR1's prefix is reachable at MR1's care-of address.

- 12 -

MR1 150 de-tunnels the packet from HA1 240 and forwards the content to the mobile router MR2 260 of the originally intended recipient, LFN2 165. Meanwhile, MR1 150 sends a Binding Update to the sender of the
5 encapsulated packet (that is, CN2 255) to inform it that MR1's prefix is reachable at MR1's care-of address.

When receiving the original packet as sent by CN2 255, MR2 260 replaces its address in the destination field of
10 the packet with the address contained in the routing header (that is, LFN2 165) and forwards the data packet to the ultimate recipient.

The inventors of the present invention have identified a
15 significant problem with the scenario depicted in FIG. 4. All subsequent packets from CN2 255 to LFN2 165 will be routed in exactly the same manner as the second data packet. That is, there will be no subsequent improvement towards route optimisation. This is more clearly shown
20 in respect of FIG. 5.

Referring now to FIG. 5, a known binding cache 500 is illustrated. The binding cache comprises a list of entries, specific to each MR in a nested mobility
25 scenario. The binding cache entries include, for example a MR3 prefix and prefix length 530, with a link 532 to a determined MR3 care-of-address 534, if one has been determined. The MR3 entry 535 is linked 536 to the next entry in the binding cache, namely that for MR2. The MR2
30 prefix and prefix length 520, includes a link 522 to a determined MR2 care-of-address 524, if one has been

- 13 -

determined. A similar arrangement and link 526 is performed to MR1, and so on.

In addition, the binding cache entries include a flag
5 entry (not shown). A 'P' flag is the "Prefix Scope
Registration" flag. When it is set, a "Home Address"
field is filled with the Mobile Network prefix (the
prefix that is advertised by the Mobile Router) and the
"Prefix Length" corresponds to the length of the Mobile
10 Network prefix.

It is specified in the document by Thierry Ernst: IETF
Internet-Draft draft-ernst-mobileip-v6-network-02.txt,
June 2001, that the Binding Cache is searched for an
15 entry corresponding to the destination address of the
packet in one pass. As a result of the search, either
nothing has been found (no entry), or the full address
has been found (128-bit match for an IPv6 address, P flag
unset), or the first bits of the destination address
20 match with a registered prefix for the registered prefix
length. In the latter case, the destination is located
in a mobile network.

Therefore, with reference to FIG. 4, when CN2 255 has to
25 send a packet to LFN2 165, CN2 255 still reviews its
binding cache and finds the entry 'MR2 260 prefix
reachable at MR2 260 Co@' 520, 524. CN2 255 does not
even consider the entry 'MR1 150 prefix reachable at MR1
Co@' 510, 514, as this would appear to have no bearing on
30 the LFN2 address. The inventors have recognised that
this deficiency results from the LFN2 165 address being
unrelated to the MR1 prefix. The fact that MR2 260 Co@

- 14 -

belongs to MR1 prefix is neither seen, nor even used by CN2 255.

Consequently, the only optimisation that the Thierry
5 Ernst proposal can support is related to the HA (HA2 250)
of the MR (MR2 260) serving the communicating MNN (LFN2
165 in the above example). This solution describes
indeed a means for having packets be sent directly from
CN2 255 to HA1 240, instead of CN2 255 to HA2 250 and
10 thereafter to HA1 240. However, if there were n
successive levels of nested mobility, this solution
provides minimal route optimisation, no more than having
a CN2 255 \rightarrow HA_{n-1} \rightarrow HA_{n-2} \rightarrow ... \rightarrow HA1 path instead of CN2
255 \rightarrow HA_n \rightarrow HA_{n-1} \rightarrow HA_{n-2} \rightarrow ... \rightarrow HA1 path. This proposal
15 is therefore still clearly inefficient, particularly in
the case of nested networks.

A need therefore arises for a mechanism, apparatus and
associated methods to support route optimisation in
20 Network mobility, especially in the case of IPv6. In
particular, a need has arisen to support route
optimisation in the case of nested mobility, wherein the
aforementioned problems are substantially alleviated.

25 **Statement of Invention**

In accordance with a first aspect of the present
invention, there is provided a method of transmitting a
data packet from a first communication node to a second
30 communication node in a mobile network, as claimed in
Claim 1.

- 15 -

In accordance with a second aspect of the present invention, there is provided a communication message, as claimed in Claim 13.

- 5 In accordance with a third aspect of the present invention, there is provided a communication message, as claimed in Claim 14.

- 10 In accordance with a fourth aspect of the present invention, there is provided a communication message, as claimed in claim 16.

- 15 In accordance with a fifth aspect of the present invention, there is provided a communication node, as claimed in Claim 17.

- 20 In accordance with a sixth aspect of the present invention, there is provided a communication node, as claimed in Claim 18.

In accordance with a seventh aspect of the present invention, there is provided a storage medium, as claimed in Claim 19.

- 25 In accordance with an eighth aspect of the present invention, there is provided a method for building an extended binding cache, as claimed in claim 20.

- 30 In accordance with a ninth aspect of the present invention, there is provided a method for constructing and sending a care-of route advertising message at a mobile network node, as claimed in Claim 21.

- 16 -

In accordance with a tenth aspect of the present invention, there is provided method for constructing and sending mobile network prefix advertising message at a mobile router, as claimed in Claim 22.

In accordance with an eleventh aspect of the present invention, there is provided a storage medium, as claimed in Claim 23.

In accordance with a twelfth aspect of the present invention, there is provided an apparatus, as claimed in Claim 24.

In accordance with a thirteenth aspect of the present invention, there is provided a communication unit, as claimed in Claim 25.

In accordance with a twelfth aspect of the present invention, there is provided a communication system, as claimed in Claim 26.

Further aspects of the present invention are as claimed in the dependent Claims.

Brief Description of the Drawings

FIG. 1 illustrates the movement of a Mobile Network in an Internet;

FIG. 2 illustrates a known packet data routing mechanism for mobile networks when applied to nested mobility;

- 17 -

FIG. 3 illustrates a known packet data routing mechanism for mobile networks when applied to nested mobility that highlights the inefficiency of the data routing;

5

FIG. 4 illustrates a known packet data routing mechanism for mobile networks when applied to nested mobility that highlights the inefficiency of an improved process of the data routing; and

10

FIG. 5 illustrates a known binding cache for routing data packets in mobile node networks.

Exemplary embodiments of the present invention will now be described, with reference to the accompanying drawings, in which:

FIG. 6 illustrates a network topology for advertising a mobile router mobility within mobile networks, in accordance with the preferred embodiment of the present invention;

FIG. 7, FIG. 8 and FIG. 9 illustrate flowcharts for an MNN router constructing and sending a care-of route advertisement, in accordance with the preferred embodiment of the present invention;

FIG. 10 illustrates a network topology for sending extended BUs to CNs in accordance with the preferred embodiment of the present invention;

- 18 -

FIG. 11 illustrates a flowchart of an MNN generating a care-of route, in accordance with embodiments of the present invention;

5 FIG. 12 and FIG. 13 illustrate flowcharts of an MNN sending extended BU messages to its CN(s) (and HAS), in accordance with the preferred embodiment of the present invention;

10 FIG. 14 and FIG. 15 illustrate flowcharts of a MNN sending an extended BU to its CN(s) (and its HAS) in accordance with the preferred embodiment of the present invention;

15 FIG. 16 illustrates a flowchart of a dynamic discovery method of a mobile network prefix for a MNN in a mobile network, in accordance with the preferred embodiment of the present invention;

20 FIG. 17, FIG. 18, and FIG. 19 illustrate flowcharts of processes of MNN routers sending a mobile network prefix advertisement message, in accordance with the preferred embodiment of the present invention;

25 FIG. 20 illustrates a flowchart of a first node sending a data packet to a second node, in accordance with the preferred embodiment of the present invention;

FIG. 21 illustrates preferred examples of data packet
30 formats sent by a first node to a second node, in accordance with embodiments of the present invention;

- 19 -

FIG. 22 and FIG. 23 illustrate a mobile network prefix solicitation message and a mobile network prefix advertisement message respectively, in accordance with the preferred embodiment of the present invention;

5

FIG. 24 and FIG. 25 illustrate a care-of route solicitation message and a care-of route advertisement message respectively, in accordance with the preferred embodiment of the present invention;

10

FIG. 26 and FIG. 27 illustrate an extended BU message sent from a LFN and a VMN respectively, in accordance with the preferred embodiment of the present invention;

15

FIG. 28 illustrates an extended binding cache in accordance with the preferred embodiment of the present invention;

20

FIG. 29 illustrates the construction of a care-of-source route from a received extended BU in accordance with the preferred embodiment of the present invention; and

25

FIG. 30 is a flowchart that illustrates a process of a first node receiving an extended BU from a second node in accordance with embodiments of the present invention.

Description of Preferred Embodiments

There is currently no standard mechanism to support Network mobility adequately, particularly in the case of IPv6 for nested mobility data networks. In particular, there is no provision or support for route optimisation.

- 20 -

This is a major issue, as nested mobility is a very realistic scenario for Mobile Router applications. In addition route optimisation is much more important with regard to movement of a Mobile Network, than for movement
5 of a single Mobile Node, since the amount of traffic handled is much greater.

The preferred embodiment of the present invention is described with respect to route optimisation in
10 transmitting at least one data packet between a corresponding node (communication unit) and a local fixed node (or vice versa). In this context, route optimisation may be viewed as "the-shortest-path" direct communication between a Mobile Network Node (MNN) and its
15 Correspondent Nodes (CNs), particularly for an MNN within a Mobile Network where a nested mobility (Mobile Networks visiting Mobile Networks) scenario exists.

It is envisaged that the corresponding node may be any
20 communication unit capable of sending a data packet across a data network (such as the Internet), for example, a web server, a PC, or a workstation running a web server such as www.yahoo.com, etc. The CN may also be any mobile data communication unit such as a general
25 packet radio system (GPRS) device or a 3rd generation (3G) cellular phone, a personal digital assistant (PDA), etc., connected to the data network through any type of access. Note that the CN may also itself be located in a mobile network.

30

As explained in the previous section, a Mobile Router (MR) must send a Binding Update (BU) message to its Home

- 21 -

Agent (HA) when moving to a Foreign Network. By sending a BU message to its HA, the MR maintains IP reach-ability for itself as well as for nodes it serves (MNNS).

5 In accordance with the preferred embodiment, no restriction is placed on the mechanism to be used by the Mobile Network to inform its Home Agent about its new location. In this regard, any known approach such as those described in [3] Thierry Ernst, Hong-Yon Lach, IETF
10 Internet-Draft draft-ernst-mobileip-v6-network-02.txt, June 2001, or [5] T. J. Kniveton, IETF Internet-Draft draft-kniveton-mobrtr-00.txt, November 2001, may be adapted and used. The adaptation of such techniques enables route optimisation to be achieved in cases of
15 nested Mobility.

Furthermore, in an enhanced embodiment of the present invention, it is shown how the approaches in [3] Thierry Ernst, IETF Internet-Draft draft-ernst-mobileip-v6-
20 network-02.txt, June 2001, and [5] may be leveraged in order to also achieve route optimisation on the path HA → MN (where MN is either a host or a router).

The preferred embodiment of the present invention
25 utilises three key concepts, which are used either singly or preferably in combination. The three key concepts are:

(i) Advertising a Mobile Router's mobility within
30 its Mobile Network.

- 22 -

(ii) Sending of one or more extended Binding Update messages by MNNs to their respective CNs. An extended Binding Update message sent by a MNN would include a "care-of route" instead of a single care-of address in vanilla Mobile IP. The care-of route is an ordered list of IP addresses that a CN will use to source route its packets to the MNN on the shortest path, thereby achieving route optimisation.

(iii) Receiving of extended BU messages by the CNs. In response to the received BU messages, source route packets addressed to MNN through the route derived from the care-of route address in order to achieve route optimisation.

15

(A) Mobile Router Advertising its mobility in the Mobile Network:

In order to realize route optimisation for MNNs in the case of a single Mobile Network, as well as multi-layered aggregated Mobile Networks (nested mobility), the inventive concepts described herein propose to make MNNs aware of the mobility of the Mobile Networks (or Mobile Routers) that they are attached to. This is in direct contrast to the approach proposed in [3] where a Mobile Router's mobility is hidden to its respective MNNs.

In order to announce it has moved to a new point of attachment, a Mobile Router (e.g. MR1) will send a "Care-of Route Advertisement" message into the Mobile Network it serves addressed to all the nodes behind it (i.e. any

- 23 -

LFNs, LMNs, VMNs). This message contains the Care-of address of the MR1 itself, as well as the care-of addresses of all the Mobile Routers above MR1 in the hierarchy of aggregated Mobile Networks, in the case of
5 nested network mobility. This list of above mobile routers' care-of addresses is preferably ordered and constructed dynamically through all the Care-of Route Advertisement sent at each level of the hierarchy of aggregated Mobile Networks.

10

Referring now to FIG. 6, a network topology 600 for advertising such a mobile router's mobility within mobile networks, is illustrated, in accordance with the preferred embodiment of the present invention. In
15 particular, the topology shown in FIG. 6, illustrates the supporting nested mobility, in accordance with the preferred embodiment of the present invention. A skilled artisan would recognise that the number of elements shown in the network of FIG. 6 are limited for clarity purposes
20 only.

The network topology shown in FIG. 6 illustrates a second Mobile Network (MR2) visiting a first Mobile Network (MR1) that is attached to a visited network. The first
25 Mobile Network includes a fixed router (LFR1) that serves its own link.

Since MR1 650 is visiting a foreign network 110, it has acquired a Care-of Address {MR1_CoA} 652. The foreign
30 network 110 is fixed, so there is no Care-of Route Advertisements forwarded in the link between the visited link 110 and MR1 650. In accordance with the preferred

- 24 -

embodiment of the present invention, MR1 650 constructs its own Care-of Route Advertisement message and includes its own Care-of Address 652 in the advertisement message.

5 This message is multicasted to all nodes within (below) its own link (MR1 Link 155) through its ingress interface. All nodes on the link will therefore receive this advertisement message. When receiving such a message, Routers (LFR, LMR, VMR) are expected to extract
10 the ordered list of Care-of addresses 652 and forward this list on to the links that they serve. In this regard, the second MR (MR2) 660 forwards the Care-of addresses 652 to its own link (MR2 link) 230. If this router is a top router of a Mobile Network not within its
15 home network (i.e. a VMR as in the case of MR2 660), it will append its own Care-of address to the list received. MR2 660 will then include the MR2 Care-of address 662 in the generation of a new ordered-list. MR2 will then generate its own Care-of Route Advertisement message to
20 be multicasted on its own link (MR2 link) 230, through its ingress interface. As a consequence, MR1_CoR is advertised in the first Mobile Network (including MR1 link 155 and LFR1 link 670) whilst the ordered-list {MR1_CoR 652, MR2_CoR 662} is advertised in the second
25 Mobile Network 230.

The process of informing CN2 655 of the optimum route to LFN2 665 via the MR2 care-of-address 662, using an extended BU message, is described later. In this manner,
30 the full route for the data packet can be determined by improved extraction, utilisation and advertisement of address data throughout the networks.

- 25 -

The process shown in FIG. 6 illustrates a practical scenario, where many nested levels may exist. Hence, a skilled artisan would appreciate that the aforementioned process may be easily generalised to an example of nested mobility involving n successive levels (n mobile routers MR_1, \dots, MR_n and n corresponding home agents HA_1, \dots, HA_n).

Referring now to FIG. 7, FIG. 8 and FIG. 9, various flowcharts 700, 800, 900 are illustrated for any router in a mobile network, such as MR VMR LFR and LMR, to construct and send a list of care-of addresses to be included in a Care-of Route advertisement (CoR_Advt) message, in accordance with the preferred embodiment of the present invention. Also, these processes are valid only for MNN that are routers (e.g. LFR, MRs).

Basically, a router in a Mobile Network will send a CoR_Advt message in response to one of three events:

- (i) Reception of a CoR_Advt on its egress interface that modifies its own CoR, as described in FIG. 7;
- (ii) Periodic sending of a CoR_Advt, as described in FIG. 8; and
- (iii) Reception of a Care-of route solicitation (CoR_Sol) message on an ingress ifc, as described in FIG. 9.

In FIG. 7, the process starts at step 710. A new CoR_Advt message is received on its egress interface, as shown in step 720, and the list of care-of addresses extracted, as described further in process A of FIG. 11. If the respective MNN's Care-of Route (MNN_CoR) has not

- 26 -

been changed in step 740, the process ends in step 780. However, if the respective MNN_CoR address has been changed in step 740, the router builds a new CoR_Advt message and includes (appends) in it the new MNN_CoR, as
5 shown in step 750. The new CoR_Advt message is then sent to all nodes through all ingress interfaces of the respective MNN, as in step 760.

In FIG. 8, a flowchart illustrates the preferred
10 operation if a periodic CoR_Advt message timer has expired in step 820. In this case, the router builds a new CoR_Advt message and includes (appends) in it to the MNN_CoR, as shown in step 850. The new CoR_Advt message is then sent to all nodes through all ingress interfaces
15 of the respective MNN, as in step 860. The CoR_Advt message timer is then re-started in step 870.

A router in a mobile network (MR, LFR, LMR, VMR) will start periodic transmission of CoR_Advt messages only
20 when its CoR becomes non-null (i.e. contains at least one address). It is envisaged that the router will terminate this periodic emission when its CoR becomes null/empty. For example, a MR will start the periodic emission when it (or a top MR) is moving out of its home network and
25 terminate the periodic transmission when it (or the top MR) is returns to its home network.

In FIG. 9, a flowchart illustrates the preferred process upon reception of a Care-of route solicitation (CoR_Sol)
30 message on an ingress ifc, as shown in step 920. In this case, the router builds a new CoR_Advt message and includes (appends) in it MNN_CoR, as shown in step 950,

- 27 -

when the (CoR_Sol) message is received on an ingress interface ifc_j. Two possible approaches are envisaged for the sending of a CoA_Advt operation in step 960. A first approach is to send the CoA_Advt message to the
5 source address of the care-of Route Solicitation message (CoR_Sol) through ifc_j. Alternatively, the CoA_Advt message may be sent to all nodes on the links through the ifc_j. Note that when a mobile router (MR) changes its location, it should send a CoR_Sol message, in order to
10 receive a CoR_Advt message in return. The MR is then able to compute its new Care-of Route (CoR). This CoR will be made of the ordered list of addresses received in the CoR_Advt message (if not empty) to which MR's new Care-of address is appended. Since this new CoR is not
15 empty, the MR will then start periodic sending of its own CoR_Advt messages.

Several approaches are envisaged for the implementation of Care-of Route Solicitation and Care-of Route
20 Advertisement messages. First, the messages may be implemented as an Internet Control Message Protocol for IPv6, IETF RFC 1885 (ICMPv6) extension, or as any new protocol on top of IP(v4 or v6), such as User Datagram Protocol, IETF RFC 768 (UDP), Transmission Control
25 Protocol, IETF RFC 793 (TCP), etc.

Secondly, a Care-of Route Advertisement message may be sent to the IPv6 all-node link-local multicast address. In this case, the message will be sent only on the local
30 link. A preferred manner would be to send the Care-of Route Advertisement to the IPv6 all-node site-local multicast address, where the site is the whole Mobile

- 28 -

Network served by this MR. This would help reduce operation on intermediate routers (LFR, LMR at home) in the Mobile Network since the Care-of Route Advertisement would then be forwarded transparently to all the links of the Mobile network. In this manner, there is no need for this intermediate router (LFR, LMR at home) to extract and copy the CoR list.

A preferred implementation of these messages would be to define them as extensions of ICMPv6 Router Solicitation and ICMPv6 Router Advertisement messages. A Care-of Route Advertisement message would be an ICMPv6 Router Advertisement message (RA) including a new "Care-of Route Advertisement" option. A Care-of Route Solicitation message would be an ICMPv6 Router Solicitation message (RS) including a new "Care-of Route Advertisement" option. The Care-of Route Advertisement option would be included in RA when a router has to announce its (non-empty)_ CoR within the mobile network. The Care-of Route Solicitation option may be included in RS by a MNN to explicitly request for a RA with the Care-of Route Advertisement option. If this option were not included in the RA, this would mean that the CoR of this router is null/empty.

Referring now to FIG. 24 and FIG. 25, a Care-of Route Solicitation message and a Care-of Route Advertisement message are illustrated respectively, in accordance with the preferred embodiment of the present invention.

The Care-of Route Solicitation (CoR_Sol) message 2400 in FIG. 24 is an ICMPv6 Router Solicitation message,

- 29 -

extended with the new "Care-of Route Solicitation"
option. It includes a single IP header 2425 that
includes an IP source address 2410 for a host and an IP
destination address 2420 indicating the all-routers IP
5 multicast address. The IP header 2425 is followed by a
router solicitation message 2430 containing the Care-of
Route Solicitation (CoR_Sol) option.

The Care-of Route Advertisement (CoR_Advt) message 2500
10 in FIG. 25 is an ICMPv6 Router Advertisement message
extended with the new "Care-of Route Advertisement"
option. It includes a single IP header 2525 that
includes an IP source address 2510 for a router's ingress
interface and an IP destination address 2520 indicating a
15 node (or all nodes on the link) that is to receive the
packet. The IP header 2525 is followed by a router
advertisement message 2530 containing the Care-of Route
Advertisement option that includes the Care-of Route
(i.e. ordered list of addresses) as advertised by the
20 router.

**B) MNNs sending extended Binding Updates to their
respective CNs:**

25

Any MNN, including a LFN, a LMN or a VMN, in a Mobile
Network is expected to send so-called extended Binding
Updates to its CNs in order to achieve route optimisation
on the path CN → MNN. In accordance with the preferred
30 embodiment of the present invention, an extended BU
message sent by a MNN would include a "care-of route"
instead of a single care-of address, as known in Mobile

- 30 -

IP. As described above, the care-of route is the ordered-list of IP addresses for the CN to derive a source route from, in order to route its packet to the MNN through the shortest path (i.e. route optimisation).

5

Referring now to FIG. 10, a network topology 1000 for sending extended BUs to CNs, is illustrated in accordance with the preferred embodiment of the present invention. For clarity purposes only, the topology follows on from that described above with respect to FIG. 6.

Each MNN in the Mobile Network derives its care-of route from the Care-of Route Advertisement messages it receives from its upper router. For instance, LFN1 675 care-of route is a single hop route equal to MR1 650's Care-of Address {MR1_CoA} 652. LFN1 675 then transmits an extended BU message 1010 indicating its care-of route to CN1 1030.

20 In contrast, LFN2 665's care-of route is a two-hop route equal to {MR1_CoA → MR2_CoA}. In this context, the LFN2 665 Care-of Route utilises MR2 CoA 662 and the upper router from MR2, namely MR1 650's CoA 652. LFN2 665 then transmits an extended BU message 1020 indicating the MR1_CoA 652 and the MR2 CoA 662 to CN2 655.

Referring now to FIG. 11, a preferred algorithm for a MNN to generate its own care-of route (MNN_CoR), in accordance with embodiments of the present invention, is described. The process (process A) starts at step 1110. When the MNN receives a new CoR_Advt message on its interface ifc_i, in step 1120, it determines whether it

30

- 31 -

is an egress interface, in step 1130. If it is not an egress interface, the CoR_Advt is ignored and the process ends in step 1190. In this case, the MNN Care-of Route is unchanged.

5

If ifc_I in an egress interface, then the MNN extracts a new CoR (new_CoR) - ordered list of IP addresses - from the CoR_advt in step 1140, and a determination is made as to whether the MNN is at home (i.e. ifc_i is attached to its home IP subnet) in step 1150. If the MNN is at home, the MNN_CoR is replaced by the new Care-of Route information in step 1180, if there is a difference between the routes in step 1170. If the MNN is not at home, the MNN care-of route is constructed by appending the MNN care-of address (obtained in the visited network) to the ordered-list of care-of addresses received in the care-of route advertisement message, in step 1160.

Note that a MNN may not maintain such a Care-of Route if it does not intend to benefit from route optimisation on the path CN→MNN, or on the path HA→MNN. However, in the preferred embodiment of the present invention, it is recommended that such Care-of Route knowledge be maintained, in order to leverage the approaches described in [3] and [5]. In this manner, it is possible to achieve route optimisation on the path HA → MN, where the MN is either a host or a router, as described later.

Note that the approach described in FIG. 11 is valid for any MNN, either a router or a host, either fixed or mobile.

- 32 -

Referring now to FIG. 12 and FIG. 13, flowcharts of an MNN sending extended BU messages to its CN(s) (and Home Agent - HA), are illustrated in accordance with the preferred embodiments of the present invention. The
5 flowcharts are again applicable to any MNN, either a host or a router, fixed or mobile.

FIG. 12 describes a flowchart 1200 for sending extended BU messages to its CN(s) (and HA) following a reception
10 of a CoR_Advt message on its egress interface, in step 1215. Upon reception of a CoR_Advt message, the MNN performs process 'A' described in FIG. 11, to generate its own CoR, as shown in step 1220. If its CoR has not changed in step 1225, the process ends, in step 1265.
15 However, if the CoR has changed in step 1225, the MNN determines all of its CNs that need to receive an updated CoR message, in step 1230. In order to send an extended BU message containing the updated CoR message, in step 1240, the MNN steps through each of the CNs, starting
20 with the first CN in step 1235. If the CN receiving the extended BU message containing the updated CoR message transmission is not the last CN identified by the MNN, in step 1245, the process moves to the next CN in step 1250 until all CNs have received the transmission.

25 A preferred, optional step is for the MNN, which is not at home, to send a message including its new Care-of Route (CoR) to its Home Agent (HA), either an Extended BU or an Extended Prefix Scope BU (if the MNN used [3] to
30 send BU to its HA), as shown in steps 1255 and 1260.

- 33 -

FIG. 13 describes an alternative flowchart 1300 for sending extended BU messages to its CN(s) (or its HA) following periodic sending of extended BUs (EBUs) to CNs (and again HA if MNN is a mobile host or router). Again, this periodic sending takes place only when the MNN's CoR is non-null following expiration of the periodic EBU timer associated to a particular CN (e.g. CNi) or MNN's HA, in step 1320. The extended BU message containing the updated CoR message is sent to CNi (or HA) in step 1330. The timer function, associated to the respective CN (e.g. CNi or HA), is then re-started in step 1340.

Note that in relation to FIG. 12 and FIG. 13, a CN may be a fixed or mobile host in the Internet as well as another MNN (from the same or different Mobile Network).

Still referring to FIG.12 and FIG. 13, if the MNN (host or router) is at home (i.e. is a LFN, a LMN at home), the MNN sends an extended Binding Update to all its CNs (and not to its HA since it is at home). If the MNN (host or router) is in a visited network (i.e. is a VMN), the MNN sends an extended Binding Update to all its CNs as well as preferably its Home Agent.

The MNN may send an extended Binding Update to its Home Agent, in case it wants to achieve route optimisation on the path HA → MNN. This approach extends the techniques proposed in [3] and [5], by sending extended BUs instead of their basic Binding Updates. In the case of [3], the Mobile Router (MR) sends a so-called prefix scope BU message to its HA. In this context, we define here an extension of this binding update message called 'extended

- 34 -

prefix-scope binding update', which includes the MR Care-of route instead of only MR care-of address. This extended prefix-scope binding update may be sent by the MR to its HA to achieve route optimisation on the path HA
5 → MR.

In the case of [5], the Mobile Router (MR) sends a basic Mobile IP binding update. In this context, we define an extension of this BU message as an '*extended Binding*
10 *Update'* (as detailed above). This includes the MR Care-of Route instead of only MR care-of address. This extended binding update may be sent by the MR to its HA to achieve route optimisation on the path HA → MR.

15 It is worth mentioning that each MNN that sends extended BU messages preferably maintains an '*extended Binding List'*, defined as an extension of the Mobile IP binding list, where Care-of Addresses are replaced with Care-of Routes.

20 Referring now to FIG. 14 and FIG. 15 illustrate flowcharts of an optimised process for an MNN sending an extended BU to its CN(s) (and its HA) for intra-mobile-network communication, in accordance with an enhanced
25 embodiment of the present invention. It is envisaged that the enhanced process improves network performance by avoiding sending "useless" extended BU messages in the case of intra-mobile-network communication. In fact, when two nodes (either fixed, or mobile at home) in the
30 same mobile network (i.e. not separated by a Mobile Router away from home) are sending packets to each other, route optimisation will be realised naturally by the

- 35 -

routing infrastructure within the mobile network. As such, there is no need for them to exchange extended BU messages between each other.

5 In FIG. 14, as an extension to the flowchart of FIG. 12, when the MNN is a host at home (i.e. is a LFH, or a LMH at home), the MNN sends an extended Binding Update to only its CNs that are not in the same mobile network; as in step 1410. Also, when the MNN is a router at home
10 (i.e. is a LFR, or a LMR at home), the MNN sends an extended Binding Update to only its CNs that are not in the same mobile network, as in step 1410. If the MNN (host or router) is in a visited network (i.e. is a VMN) in step 1420, the MNN sends an extended Binding Update to
15 only CNs that are not in this visited mobile network, in step 1240. For CNs that are in the visited mobile network, the MNN may send an extended Binding Update, or preferably send a modified extended Binding Update where only one address is specified in the care-of route, i.e.
20 only the MNN's own care-of address, in step 1430.

The mobile MNN (host or router) preferably sends also a Binding Update to its Home Agent in step 1260 and 1450. Two cases are considered below, in this regard. When the
25 MNN has moved out of its home mobile network, the MNN preferably sends an extended Binding Update to its HA. When the MNN is in a foreign IP-subnet within its home mobile network, the MNN may send an extended Binding Update. Alternatively, the MNN may send a modified
30 extended Binding Update where only one address is specified in the care-of route, i.e. its own Care-of Address.

- 36 -

Again, it is envisaged that the MNN may send an extended BU message to its Home Agent, in the case where the MNN wants to achieve route optimisation on the path HA → MNN, by sending extended BUs in the operation of [3] or extended prefix scope BUs in the operation of [5] as described above.

A similar enhancement to FIG. 13 is described in FIG. 15, with the transmission of extended BU messages containing only the MNN_CoA to respective CNs that are operating in the same Mobile Network as the MNN, as shown in steps 1510, 1520, 1530.

Referring now to FIG. 26 and FIG. 27 an extended BU message sent from a LFN and a VMN are illustrated respectively, in accordance with the preferred embodiment of the present invention.

The extended BU message sent from a LFN message 2600 in FIG. 26 includes a single IP header 2625 that includes as IP source address 2610 the LFN and as IP destination address 2620 for the CN's address. The IP header 2625 is followed by a Mobile IPv6 BU message 2630 containing the Care-of Route mobility option incorporating the LFN's Care-of Route.

The extended BU message sent from a VMN message 2700 in FIG. 27 includes a single IP header 2725 that includes an IP source address 2710 of the VMN care-of address and an IP destination address 2720 for the CN. The IP header 2725 is followed by a Mobile IPv6 BU message 2730

- 37 -

containing the Care-of Route mobility option
incorporating the VMN's Care-of Route.

Referring now to FIG. 16 a flowchart 1600 of a dynamic
5 discovery method of a mobile network prefix for any MNN
in a mobile network is illustrated, in accordance with
the enhanced embodiment of the present invention
described above. In this context, a node such as a MNN
is able to know whether a second node, say its CN, is in
10 the same current Mobile Network as itself when sending an
EBU to say the CN, as described above.

The process (process B) starts at step 1610. When the
MNN receives a new Mobile Network Prefix Advertisement
15 message (Mobile_Network_Prefix_Advt) on its interface
ifc_i, in step 1620, it determines whether it is an
egress interface, in step 1630. If it is not an egress
interface, the Mobile_Network_Prefix_Advt is ignored and
the process ends in step 1670. In this case, the MNN's
20 mobile network prefix (MNN_Mobile_Netowrk_Prefix) is
unchanged. If ifc_i is an egress interface, then the MNN
extracts a new Mobile Network Prefix
(new_Mobile_Network_Prefix) and its prefix length
(new_Mobile_Network_Prefix_Length) from the
25 Mobile_Network_Prefix_Advt in step 1640, and a
determination is made as to whether the MNN's home
address in ifc_i matches New_Mobile_Network_Prefix on the
first New_Mobile_Network_Prefix_Length bits, in step
1650. If it matches, the MNN_Mobile_Network_Prefix and
30 MNN_Mobile_Netowrk_Prefix_Length are replaced by the new
information in step 1650, and the process ends in step

- 38 -

1670. If there is not a match, the process ends directly in step 1670.

Note that the approach described in FIG. 16 is valid for
5 any MNN, either a router or a host, either fixed or mobile.

Referring now to FIG. 17, FIG. 18 and FIG. 19, various flowcharts 1700, 1800, 1900 are illustrated for a fixed
10 router in a mobile network (i.e. LFR) to construct and send its own Mobile Network Prefix Advertisement (Mobile_Network_Prefix_Advt) messages, in accordance with the enhanced embodiment of the present invention.

15 Basically, a fixed router in a Mobile Network will send a Mobile_Network_Prefix_Advt message in response to one of three events:

(i) Reception of a Mobile_Network_Prefix_Advt on an
20 egress interface that modifies its own MNN_Mobile_Network_Prefix, as described in FIG. 17;

(ii) Periodic sending of a
Mobile_Network_Prefix_Advt, as described in FIG. 18; and

25

(iii) Reception of a Mobile Network Prefix Solicitation message (Mobile_Network_Prefix_Sol) on an ingress ifc, as described in FIG. 19.

30 In FIG. 17, the process starts at step 1705. A new Mobile_Network_Prefix_Advt message is received on an egress interface, as shown in step 1710, and the new

- 39 -

Mobile Network Prefix (New_Mobile_Network_Prefix) and its length are extracted, as described further in process B of FIG. 16. If the respective MNN_Mobile_Network_Prefix and MNN_Mobile_Network_Prefix_Length have not been

5 changed in step 1720, the process ends in step 1735. However, if the respective MNN_Mobile_Network_Prefix and/or MNN_Mobile_Network_Prefix_Length have been changed in step 1720, the router builds a new

Mobile_Network_Prefix_Advt message and includes (appends)

10 in it the new MNN_Mobile_Network_Prefix and MNN_Mobile_Network_Prefix_Length, as shown in step 1725. The new Mobile_Network_Prefix_Advt message is then sent to all nodes through all ingress interfaces of the respective MNN, as in step 1730. The flowchart described

15 in FIG. 17 only applies to a fixed router in a mobile network.

In FIG. 18, a flowchart illustrates the preferred operation if a periodic Mobile_Network_Prefix_Advt

20 message timer has expired in step 1810. In this case, the router builds a new Mobile_Network_Prefix_Advt message and includes (appends) in it the

MNN_Mobile_Network_Prefix and

MNN_Mobile_Network_Prefix_Length, as shown in step 1815.

25 The new Mobile_Network_Prefix_Advt message is then sent to all nodes through all ingress interfaces of the respective MNN (fixed router), as in step 1820. The Mobile_Network_Prefix_Advt message timer is then re-started in step 1825. The flowchart of FIG. 18 applies

30 to any router in a mobile network (MR, LMR, LFR, VMR).

- 40 -

A router in a mobile network(MR, LFR, LMR, VMR) will start periodic transmission of Mobile_Network_Prefix_Advt messages only when its CoR becomes non-null (i.e. contains at least one address). It is envisaged that the
5 router will terminate this periodic emission when its CoR becomes null/empty. For example, a MR will start the periodic emission when it (or a top MR) is moving out of its home network and terminate the periodic transmission when it (or the top MR) has returned to its home network.

10

In FIG. 19, a flowchart illustrates the preferred process upon reception of a Mobile Network Prefix Solicitation message (Mobile_Network_Prefix_Sol) on an ingress ifc, as shown in step 1910. In this case, the router builds a
15 new Mobile_Network_Prefix_Advt message and includes (appends) in it MNN_Mobile_Network_Prefix and MNN_Mobile_Network_Prefix_Length, as shown in step 1920, when the (Mobile_Network_Prefix_Sol) message is received on an ingress interface ifc_j. Two possible approaches
20 are envisaged for the sending of a Mobile_Network_Prefix_Advt operation in step 1925. A first approach is to send the message to the source address of the Mobile Network Prefix Solicitation message (Mobile_Network_Prefix_Sol) through ifc_j.
25 Alternatively, the Mobile Network Prefix Advertisement message may be sent to all nodes on the links through the ifc_j. The flowchart of FIG. 19 applies to any router in a mobile network (MR, LMR, LFR, VMR).

30 Note that when a mobile router (MR) changes its location (and thus get a non-null CoR), it should send a Mobile_Network_Prefix_Advt on its ingress interfaces.

- 41 -

This mobile router may know this prefix from its internal configuration. On the other hand, fixed routers in the mobile network (i.e. LFR) may not a-priori know the mobile network prefix they belong to and can discovery it
5 dynamically thanks to the process described in FIG. 17. Such fixed routers may of course send a Mobile_Network_Prefix_Sol message, in order to receive a Mobile_Network_Prefix_Advt message in return.

10 Several approaches are envisaged for the implementation of Mobile Network Prefix Solicitation and Mobile Network Prefix Advertisement messages. First, the messages may be implemented as ICMPv6 extension, or as any new protocol on top of IP, UDP, TCP, etc. ICMP, UDP and TCP
15 are protocols that build upon IP(v4 or v6):

- ICMPv6: Internet Control Message Protocol for IPv6, IETF RFC 1885
- UDP: User Datagram Protocol, IETF RFC 768
- 20 - TCP: Transmission Control Protocol, IETF RFC 793

Secondly, a Network Prefix Advertisement message may be sent to the IPv6 all-node link-local multicast address. In this case, the message will be sent only on the local
25 link. A preferred manner would be to send the Care-of Route Advertisement to the IPv6 all-node site-local multicast address, where the site is the whole Mobile Network served by this MR. This would help reduce operation on intermediate routers (i.e. LFR) in the
30 Mobile Network since this message would then be forwarded transparently to all the links of the Mobile network. In this manner, there is no need for this intermediate

- 42 -

router (i.e. LFR) to extract and copy the Mobile Network Prefix.

A preferred implementation of these messages would be to
5 define them as extensions of ICMPv6 Router Solicitation
and ICMPv6 Router Advertisement messages. A Mobile
Network Prefix Advertisement message would be an ICMPv6
Router Advertisement message (RA) including a new "Mobile
Network Prefix Advertisement" option. A Mobile Network
10 Prefix Solicitation message would be an ICMPv6 Router
Solicitation message (RS) including a new "Mobile Network
Prefix Advertisement" option. The Mobile Network Prefix
Advertisement option would be included in RA when a
router has to announce the mobile network prefix within
15 the mobile network. The Mobile Network Prefix
Solicitation option may be included in RS by a MNN to
explicitly request for a RA with the Mobile Network
Prefix Advertisement option.

20 Referring now to FIG. 22 and FIG. 23, a mobile network
prefix solicitation message and a mobile network prefix
advertisement message are illustrated respectively, in
accordance with the preferred embodiment of the present
invention.

25 The mobile network prefix solicitation message 2200 in
FIG. 22 is an ICMPv6 Router Solicitation extended with
the new Mobile Network Prefix Solicitation option. It
includes a single IP header 2225 that includes an IP
30 source address 2210 for a host and an IP destination
address 2220 indicating the all-routers IP multicast
address. The IP header 2225 is followed by the router

- 43 -

solicitation message 2230 containing the mobile network prefix solicitation option proposed in this document.

The mobile network prefix advertisement message 2300 in
5 FIG. 23 is an ICMPv6 Router Advertisement message
extended with the new mobile network prefix advertisement
option. It includes a single IP header 2325 that
includes an IP source address 2310 for a router's ingress
interface and an IP destination address 2320 indicating a
10 node (or all nodes on the link) that is to receive the
packet. The IP header 2325 is followed by the router
advertisement message 2330 containing the mobile network
prefix advertisement option that includes the mobile
network prefix and the mobile network prefix length
15 advertised by the router.

**C) Correspondent Nodes and Home Agent receiving extended
Binding Update:**

20 Referring now to FIG. 20, a flowchart 2000 illustrates a
process for a first node to send a data packet to a
second node based on the parsing of its Extended Binding
Cache (EBC), in accordance with the preferred embodiments
25 of the present invention. A first node, for example a CN
or HA, receives an indication that a data packet is to be
sent to a second node, for example a MNN, as shown in
step 2020. The CN searches for the MNN's address within
an extended binding cache (EBC) that stores the Care-of
30 Source routes derived from care-of route information
received in extended BUs, as in step 2030.

- 44 -

If the MNN's address is found, in step 2040, the CN (or HA) is able to source route the packet, through the Care-of Source route found in the EBC, to the MNN's address achieving route optimisation, as in step 2050.

- 5 Otherwise, the CN sends the data packet to the MNN directly to MNN's Home Address using the known technique, as shown in step 2060.

Referring now to FIG. 21, preferred examples of data
10 packet formats, sent by a first node, say a CN, to a second node, say a MNN, are illustrated in accordance with the preferred embodiments of the present invention.

A first format of a data packet 2100 includes a single IP
15 header 2110 that includes an IP source address 2112 and an IP destination address 2114 equal to the first IP address in CN's Care-of Source Route to MNN. The single IP header 2110 is followed by a routing header 2120 containing the (m-1) other addresses (from CN's Care-of
20 Source route to MNN) in the route to the second node. The data payload 2130 follows the routing header.

A second format of a data packet 2150 includes 'm'
consecutive IP headers each including respective IP
25 source addresses 2112, 2152, 2162, 2172 and respective IP destination addresses 2114, 2154, 2164, 2174. The 'm' consecutive IP headers do not therefore need a separate routing header, so the data payload 2180 follows the IP headers. Each of the consecutive IP headers has its IP
30 destination address set equal to one of the IP addresses of CN's Care-of Source route to MNN. The first header contains the first address of the Care-of Source Route

- 45 -

(CoSR), the second header contains the second address, and so on to the last header. The final IP header 2172, 2174, together with the data payload 2180, constitutes the original packet 2190 to be sent from the first node
5 to the second node.

Each node (CN, MNN, HA) that is expected to receive extended Binding Updates is to maintain an extended Binding Cache (EBC). The EBC is herein defined as an
10 extension of Mobile IP binding cache where care-of addresses are replaced with "care-of source routes" (CoSR) derived from the care-of routes (CoR) received in the extended Binding Updates. It is worth mentioning that the care-of source route listed in the extended
15 binding cache may be slightly different (i.e. shorter) from the care-of route received in the extended Binding Update.

Referring now to FIG. 28 an extended binding cache 2800
20 is illustrated in accordance with the preferred embodiment of the present invention.

The extended binding cache 2800 includes a list of entries 2810, 2840, 2870, each one specific to a home
25 address of a MNN. The extended binding cache entry 2810 includes, for example, a first home address 2815, with a link 2818 to a first Care-of Source route 2820, if one has been determined. The first Care-of Source route 2820 includes respective linked Care-of Source route addresses
30 2825, 2830, and 2835 to identify the route to the first home address. The first entry 2810 is linked 2837 to a second entry 2840 having a link 2848 from a second home

- 46 -

address 2845 to a second Care-of Source route 2850, if one has been determined. The second Care-of Source route 2850 includes respective linked Care-of Source route addresses 2855, 2860, and 2865 to identify the route to
5 the second home address. A similar arrangement and link 2867 is performed to a third entry 2870, and so on.

It is within the contemplation of the invention that the extended BU may also include entries for a mobile router
10 (MR) prefix and prefix length. This is particularly useful when an MR sends an extended prefix-scope BU to its HA or CNs, as an extension to [3], where the CoR replaces the CoA in the prefix-scope binding update. In this case, the Home Address field is replaced with a
15 "prefix and prefix length" field in the EBC of the respective HA or CNs.

Referring now to FIG. 29, the construction of a care-of-source route 2900 from a received extended BU is
20 illustrated, in accordance with the preferred embodiment of the present invention. A first node Care-of Route (N1_CoR) contains an ordered list of Care-of Route addresses (N1_CoR[1], N1_CoR[2], ...) 2910. When the first node N1 receives an extended BU message from a second
25 node N2 containing N2's Care-of Route (N2_CoR), the first node compares the two ordered list of Care-of Route addresses, to determine when there is a difference between them 2940. The address from N2_CoR that was determined as being different, and all subsequent
30 addresses from N2_CoR, is then used to generate a new Care-of Source Route 2930 for data packet transmissions

- 47 -

from N1 to N2. This process is further described with respect to the flowchart of FIG. 30.

Referring now to FIG. 30, a flowchart 3000 a process of
5 determining a Care-of Source Route to be included in an extended Binding Cache is illustrated, in accordance with preferred embodiments of the present invention. A first node (N1) receives an extended Binding Update from a second node (N2) and needs to determine the care-of
10 source route to be included in the extended Binding Cache for this entry (N2).

The process starts at step 3002. When a first node (N1, which may be itself a MNN at home or in a foreign
15 network, or any host in the topology) receives an EBU containing a new Care-of Route for a second node (N2), in step 3004, N1 determines whether this N2 Care-of Route (received in the EBU) is empty in step 3006.

20 If the new Care-of Route for N2 received in the EBU is empty, in step 3006, N1 searches through its extended Binding Cache to check whether there is an entry (i.e. a care-of source route) for this second node in the EBC, as shown in step 3008. If no entry exists in the EBC, the
25 process ends in step 3046. In this case, no entry in N1's extended binding cache is needed for N2. When addressing packet to N2's home address, the data packet from N1 will be directly on the shortest path (route optimisation is achieved without needing source routing).
30 However, if an entry exists in the EBC in step 3008, the second node entry is removed (i.e. N1_CoSR(to N2)) in step 3010, before the process ends at step 3046.

- 48 -

If the new N2 Care-of Route received in the EBU is not empty, in step 3006, a determination is made in step 3012 as to whether a Care-of Route entry is available for the first node. If a Care-of Route is not available for the first node in step 3012, a first node (N1) Care-of Source Route to the second node (N2) is generated to be equal to the second node Care-of Route (N1_CoSR(to N2) := N2_CoR), as shown in step 3014. N1 then searches through its extended Binding Cache to check whether there is an entry (i.e. a care-of source route) for this second node in the EBC, as shown in step 3016. If no entry exists in the EBC, an entry for N2 is added to the EBC (N1_CoSR(to N2)), in step 3020 and the process ends in step 3046. If an entry exists in the EBC in step 3016, the second node entry is updated (i.e. N1_CoSR(to N2)) in step 3018, before the process ends at step 3046.

If the first node has its own Care-of Route in step 3012, all addresses in the two nodes care-of routes (N1 and N2) are searched, starting with setting a counter (i=0) in step 3022. N1 compares, one by one, the IP addresses in its own care-of route (N1_CoR) with the IP addresses listed in the care-of route of the extended Binding Update received from node N2, in step 3024.

It starts with the first addresses N1_CoR(1) and N2_CoR(1), and continues thereafter. If a match for a particular address of the first and second nodes care-of routes is found, in step 3026, a determination is made as to whether the second node Care-of Route address is the last address, in step 3028. If it is the last address,

- 49 -

the complete Care-of Route of N2 has been searched, in step 3030, and the process moves to step 3008 as described above. If the second node Care-of Route address is not the last address, in step 3028, a

5 determination is made as to whether the first node Care-of Route address is the last address for the first node, in step 3032. If the first node Care-of Route address is not the last address, in step 3032, the counter is incremented in step 3034, and the next addresses searched

10 in step 3024. If the first node Care-of Route address is the last address, in step 3028, or no match was found in step 3026, the searching process stops in step 3036 and the care-of source route to N2, from N1, is set in 3038.

15 In 3038, the care-of source route to N2, from N1, is set to be equal to ordered list of addresses part of N2's care-of route starting from the i^{th} address to the last one. This i^{th} address being the one where the loop in N2's care-of route addresses stopped in step 3036. That

20 is: $N1_CoSR(\text{to } N2) = \{N2_CoR(i) \rightarrow N2_CoR(i+1) \rightarrow \dots \rightarrow N2_CoR(n-1) \rightarrow N2_CoR(n)\}$. N1 then searches through its extended Binding Cache to check whether there is an entry (i.e. a care-of source route) for this second node in the EBC, as shown in step 3040.

25 If no entry exists in the EBC, an entry for N2 is added to the EBC ($N1_CoSR(\text{to } N2)$), in step 3042 and the process ends in step 3046. If an entry exists in the EBC in step 3040, the second node entry is updated (i.e. $N1_CoSR(\text{to } N2)$) in N1's EBC in step 3044, before the process ends at

30 step 3046.

- 50 -

If N1 has no care-of route (NR_CoR is empty or does not exist) in step 3012, then the care-of source route to N2, from N1, is set to be equal to N2's care-of route in step 3014. That is: $N1_CoSR(to\ N2) = N2_CoR$. N1 then searches through its extended Binding Cache to check whether there is an entry (i.e. a care-of source route) for this second node in the EBC, as shown in step 3016. If no entry exists in the EBC, an entry for N2 is added to the EBC ($N1_CoSR(to\ N2)$), in step 3020 and the process ends in step 3046. If an entry exists in the EBC in step 3016, the second node entry is updated (i.e. $N1_CoSR(to\ N2)$) in N1's EBC in step 3018, before the process ends at step 3046.

N1 should then source route the packet via the care-of source route to N2 in order to achieve route optimisation, as described with reference to FIG. 21 and FIG. 22. This can be implemented in several ways, as described previously. A first way is to use the IPv6 Routing Header. In this case, the first address in the care-of source route is set as the destination address in the IPv6 header and the remaining addresses of the care-of source route (CoSR) are set in the Routing Header in the same order. A second way is for the first node to build 'n' level of encapsulation of the data packet to be sent, where 'n' is the number of addresses in the care-of source route. The encapsulation #k will have the address #k, of the ordered-list of addresses in the care-of source route, as the destination address.

- 51 -

Furthermore, the various steps described above need not necessarily be performed in the order they have been described. A skilled artisan would recognise that an alternative order may be used, where benefit can still be gained in the route optimisation process.

It is to be appreciated that the arrangement and specific details of interfaces, address types, routers etc. in the above embodiments are merely examples, and the invention is not limited to these examples. The invention should be viewed as capable of being applied to other aspects of the Internet or other types of data networks or protocols and subnets thereof. Furthermore, the invention may also be applied to networks other than the Internet, when such networks have subnets and access networks corresponding to those described above for the case of the Internet.

The invention, or at least embodiments thereof, tend to provide the following advantages, singly or in combination:

(i) A data packet may be transmitted much more efficiently, using this improved route optimisation technique.

(ii) Improved privacy of data packet transmissions, as each node in the topology is responsible in sending their own Binding Updates to their CNs and Home Agent. Hence, each node is able to decide whether it wishes to disclose its current location by sending a Binding Update in order to perform route

- 52 -

optimisation. This is a clear advantage over [3] where privacy is not allowed.

(iii) The security solutions defined by the IETF, to provide "address ownership" authorisation (e.g. Return Route-ability), are still applicable in the context of the present invention. This is again a clear advantage over [3], which requires a new security mechanism to be introduced.

10

(iv) A secured exchange of new messages (e.g. "Care-of Route Advertisement") introduced in the present invention can be very easily realised, either through shared secret (in the case of nodes belonging to the same organisation) or through a new IETF protocol such as PANA in the case of a Mobile Node (host or router) visiting a foreign network.

(v) A solution for efficient data routing in IPv6, IPv4 or similar data network protocol has been achieved, particularly for systems supporting nested network mobility.

Whilst the specific and preferred implementations of the embodiments of the present invention are described above, it is clear that one skilled in the art could readily apply variations and modifications of such inventive concepts.

Thus, a mechanism, apparatus and associated methods to support route optimisation in Network mobility, especially in the case of IPv6, have been described,

- 53 -

whereby the disadvantages associated with known mechanisms, apparatus and associated methods have been substantially alleviated. In particular, a mechanism, apparatus and associated methods to support route
5 optimisation in the case of nested mobility, have been described.